Effective Scheduling Algorithms for I/O Blocking with a Multi-Frame Task Model*

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SUMMARY A task that suspends itself to wait for an I/O completion or to wait for an event from another node in distributed environments is called an I/O blocking task. Conventional hard real-time scheduling theories use framework of rate monotonic analysis (RMA) to schedule such I/O blocking tasks. However, most of them are pessimistic. In this paper, we propose effective algorithms that can schedule a task set which has I/O blocking tasks under dynamic priority assignment. We present a new critical instant theorem for the multi-frame task set under dynamic priority assignment. The schedulability is analyzed under the new critical instant theorem. For the schedulability analysis, this paper presents saturation summation which is used to calculate the maximum interference function (MIF). With saturation summation, the schedulability of a task set having I/O blocking tasks can be analyzed more accurately. We propose an algorithm which is called Frame Laxity Monotonic Scheduling (FLMS). A genetic algorithm (GA) is also applied. From our experiments, we can conclude that FLMS can significantly reduce the calculation time, and GA can improve task schedulability ratio more than is possible with FLMS.

key words: I/O blocking, multi-frame task model, schedulability analysis, laxity, genetic algorithm

1. Introduction

A task that suspends itself to wait for an I/O completion or to wait for an event from another node in distributed environments is called an I/O blocking task. With the conventional framework of the rate monotonic analysis (RMA), an I/O blocking task cannot be scheduled effectively. The reason for this is that the critical instant theorem does not follow its original form when a task has I/O blocking, and various analysis techniques derived from the theorem cannot be accordingly applied [1].

Takada [1] has proposed an approach using multi-frame task model to schedule a set of tasks including I/O blocking [2]. According to this approach, a task having I/O blocking is divided into two frames \( f_1 \) and \( f_2 \), whose maximum execution times are \( C_1 \) and \( C_2 \), respectively, one frame applicable before the blocking and the other after the blocking. The whole task period is denoted as \( T \), and the periods of \( f_1 \) and \( f_2 \) are denoted as \( T_1 \) and \( T_2 \), respectively, where \( T = T_1 + T_2 \). \( f_1 \) is started periodically within the period of \( T_1 \) and an I/O operation is requested after \( f_1 \) is completed.

Then, \( f_2 \) is started immediately after the I/O operation is completed. Here, we assume that periods \( T_1 \) and \( T_2 \) are equal to deadlines of \( f_1 \) and \( f_2 \), respectively.

In [1], a method is considered by which the deadline of a frame can be allocated so that the maximum schedulability of lower priority frames can be achieved. If subtraction of the execution time from the deadline can be described as the slack of a frame, a simple allocation method of deadlines is to make the slacks of the two frames equal. In other words, \( T_1 - C_1 = T_2 - C_2 \). In addition to the \( T = T_1 + T_2 \) formula, \( T_1 \) can be calculated by the formula \( T_1 = (T + C_1 - C_2)/2 \). This method is described as an equal slack allocation method.

Although this method is not the optimal method, compared to conventional approaches, it does provide a useful way of improving the schedulability of I/O blocking tasks.

Endo et al. have discussed the critical instant theorem for generalized multi-frame tasks (GMF) with the dynamic priority assignment [3]. By use of the maximum interference function (MIF) [4], they have calculated the interference time caused by high priority tasks within the duration time interval \( t \). Based on this theorem, they discussed the necessary and sufficient condition for the schedulability of a set of GMF tasks under dynamic priority assignment. As described later, our algorithms presented in this paper are based on their work.

Iiyama et al. [5] have applied the GMF model to engine management systems. According to their work, the degree of schedulability of a frame is captured by summation of the maximum total interference times requested by all higher priority frames within the duration of time \( t \). This summation can be calculated by “adding up” the MIFs of all higher priority frames. However, the time which obtained by normal arithmetic addition is longer than actual interference time, so they have proposed saturation summation for greater precision. We also employ saturation summation to improve the calculation of MIF in our algorithms.

As above described, scheduling of a set of tasks which has I/O blocking, allocation of the deadline and assignment of priorities to the frames simultaneously, remain as an open issue. Specifically, no approach exists to assign arbitrary priorities within a general multi-frame model. We propose an algorithm which is called Frame Laxity Monotonic Scheduling (FLMS). A genetic algorithm is also applied. As a result of our experiments, we have concluded that the genetic algorithm we proposed significantly improve task schedulability ratio more than is possible with FLMS.

The remainder of this paper is organized as follows: In
In this section, we give a brief description of the multi-frame task model under dynamic priority assignments, then give a simple example to show difference priorities can improve schedulability of a task set which has I/O blocking tasks.

2.1 The Multi-Frame Task Model

The multi-frame task model is effective for tasks whose execution times are periodically changed according to a specified pattern. Under this model, each period with the Liu and Layland (L&L) task model is called a frame, and the task maximum execution time in each frame is changed with a larger cycle. In other words, a set of tasks is divided into some frames with the same period, and the execution times are periodically changed according to a specified pattern. Under this model, each period with the Liu and Layland (L&L) task model is called a frame, and the task maximum execution time in each frame which it belongs to. A task is set at the beginning of the frame, and the deadline is the end of the frame. The generalized multi-frame task model is an extension of the multi-frame task model in the following two directions: (1) different frames may have different periods and (2) the deadline for a frame may be different from the end of the frame [6]. From now on, we will describe the generalized multi-frame task model simply as a multi-frame task model.

In addition to priorities for frames under multi-frame model, a multi-frame task \(T_i\) consisting of \(N_i\) frames is characterized by a sequence of 4-tuples \(((C^0_i, D^0_i, P^0_i, L^0_i), \ldots, (C^{N_i-1}_i, D^{N_i-1}_i, P^{N_i-1}_i, L^{N_i-1}_i))\), where \(C^j_i\), \(D^j_i\), \(P^j_i\), and \(L^j_i\) represent the maximum execution time of the \(j\)th frame of \(T_i\), its relative deadline, its priority and its length, which is its minimum separation time between the arrival time of the \(j\)th frame and the following frame. In order to maintain the frame separation properties, we assume that \(D^j_i \leq L^j_i\), holds for all \(0 \leq j \leq N_i - 1\) [4].

An I/O blocking task is shown in Fig. 1. The I/O blocking task is divided into frame1 which is before the blocking and frame2 which is after the blocking. Frame1 and frame2 are expressed as \((C1, D1, P1, L1)\) and \((C2, D2, P2, L2)\), respectively. \(D\) is the total deadline of the I/O blocking task. In this paper, we assume that \(D\) is equal to the period of the I/O blocking task. In consideration of the I/O blocking task constraints, three equations are applied as follows.

\[
L1 = D1 + B \\
D = L1 + D2 \\
P = L1 + L2
\]

The focus of our research is to find an optimal scheduling method by adjustment of priorities and deadlines of frames. To a set of period tasks according to the L&L task model, the Deadline Monotonic (DM) [7] priority assignment is optimal, meaning that if any static priority scheduling algorithm can schedule a set of tasks with deadlines that are unequal to their periods, then DM will also schedule the task set. However, DM is not the optimal scheduling method within multi-frame task model. To demonstrate this, we show an example below.

2.2 An Illustrative Example

A multi-frame task \(\tau_m\) consists of two frames \(((3, 3, P1, 3), (2, 5, P2, 5))\). There is another period task \(\tau\), the execution time, the deadline, the relative priority and the period of task \(\tau\) are 3, 6, P3 and 8, respectively. According to DM, the two frames of task \(\tau_m\) are allocated higher priority and task \(\tau\) with lower priority. Used with this kind of priority assignment, the task set is not schedulable, as shown in Fig. 2 (a).

However, if we assign \(\tau_{m,0}\) the highest priority, \(\tau\) middle priority, \(\tau_{m,1}\) the lowest priority respectively, the task set is schedulable, as shown in Fig. 2 (b).

This illustrates that DM can not effectively schedule a GMF task set. Since the release time of frames is generated by dynamic priority assignment, the problem becomes NP-hard [8]. On the other hand, an I/O blocking task allows preemption tends to make the problems easier. We consider that dynamic priority assignment for frames can improve the schedulability of a set of tasks.

2.3 Related Work

The I/O performance has been considered in some applications such as distributed systems, multiprocessor or...
multiple-disk systems [8]–[11] by other approaches. A.L. Narasimha Reddy and Prithviraj Banerjee have been improved the I/O performance by building an I/O subsystem with multiple disks [9]. Specifically, they considered disk synchronization, data declustering/disk striping, and a combination of both these approaches. The effects of block size and other parameters of the system have been evaluated. In [10], several design options in designing an I/O system have been studied. Trace driven simulations have been used to study a disk system with nonvolatile cache. It was shown that decoupling the cache block size and fetch size yielded significant performance benefits. Also, A.L. Narasimha Reddy [11] studied the I/O behavior of some scientific applications, a subset of perfect benchmarks, executing on a multiprocessor. The aim of this study is to explore the various patterns of I/O access of large scientific applications and to understand the impact of this observed behavior on the I/O subsystem architecture.

In distributed systems, the lack of global information about data transfer between clients and servers make implementation of parallel I/O a challenging task. J. Wu et al. have proposed two distributed algorithms for scheduling data transfer in parallel I/O with non-uniform data sizes [12]. Their simulations show that two algorithms are suitable for parallel I/O with data transfer traffic.

3. Schedulability Analysis

3.1 Critical Instant Theorem

With L&L task model, a critical instant of any task occurs whenever the task is requested simultaneously with requests for all higher priority tasks [13]. In case of a GMF task set which has non-AM (Accumulative Monotonic) tasks, a critical instant for the frame cannot be determined independently of the execution time of the higher priority frames. For example, consider an I/O blocking task \( \tau_m(f_1, f_2) = (3, 3, P_1, 3), (1, 5, P_2, 5) \) and a task \( \tau = (2, 2, P_3, 5) \). Assuming \( P_1 > P_3 > P_2 \), if we consider a critical instant of the task set with multi-frame model as above, a critical instant of the lowest priority frame \( f_2 \) occurs when \( f_2 \) is requested simultaneously with requests for \( \tau \). As shown in Fig. 3 (a), the response time of \( f_2 \) is 3.

However, the worst case response time of \( f_2 \) is as shown in Fig. 3 (b), and a critical instant of \( f_2 \) defined as above is not the worst case response time of \( f_2 \). Even if \( f_1 \) is executed at time 0, it can be postponed by the length of \( f_2 \). Based on the preemptive static priority scheduling theorem, we give a brief description of the new critical instant theorem for a GMF model [4].

**Definition 1 (Critical Instant Candidates):** The critical instant candidates of a GMF task are defined as the instants at which a frame of the task is requested simultaneously with any one of the frames with higher priority. Their succeeding frames are requested with their minimum separations, and each frame is executed for a maximum execution time.

On the basis of such a definition, we can describe the critical instant theorem for a GMF task as follows.

**Theorem 1** (Critical Instant Theorem): A critical instant of a frame of a GMF task occurs at one of the critical instant candidates for the task. In other words, one of the critical instant candidates is a critical instant of the frame.

**Proof.** All frames which are in a frame set \( U \) have their priorities. We assume that when frame \( F \) with priority \( i \) of the task \( T \) is released at time \( t \) and completes at time \( t_{end} \), the worst case response time of frame \( F \) is caused. In this situation, a frame is sure to be executed with higher priority than the priority of \( F \) between time \( t \) and \( t_{end} \). The time at which there is no frame with higher priority has been executed before time \( t \) is assumed to be \( t_0 \). Because any frame is not executed at time 0, \( t_0 \) is sure to exist. It is assumed that \( F \) is the first arrival of \( T \) after \( t_0 \). Even if we change the arrival time \( t \) of frame \( F \) to \( t_0 \), it still completes at time \( t \) since there is at least one frame with priority that is higher than \( i \) is executed.

Assuming that for each frame with priority that is higher than \( i \), we move its first arrival time after \( t_0 \) to \( t_0 \). In this case, the completion time of frame \( F \) might be delayed to \( t_{end'} \). Moreover, their succeeding frames are requested early with their minimum separations, and each frame is executed for a maximum execution time. The completion time of frame \( F \) is delayed further to time \( t_{end''} \). Since time \( t \) is one of critical instant candidates of \( T \), the response time of frame \( F \) is longer than the first situation.

When we assume that if the worst case response time of frame \( F \) is caused in the first situation, the worst case response time of frame \( F \) is caused too when the critical instant candidates for frame \( F \) are requested.

Next, assuming that frame \( F \) is not the first arrival of \( T \) after \( t_0 \), we prove it as follows. If frame \( F \) with priority \( i \) is the \( k \)th frame of the arrival of \( T \) after \( t_0 \), we assume that the worst case response time of frame \( F \) is caused, and it completes at time \( t_{end} \). As above proof, the time which a frame with higher priority has not been executed before \( t \) is assumed to be \( t_0 \). It is sure that \( t_0 \) exists. We assume that frame \( F' \) with higher priority \( j \) is the first arrival of \( T \).
after \( t_0 \) \((j \geq i)\). The arrival time of \( F' \) is \( t' \). It completes at time \( t_{\text{end}'} \). Even if we change at the arrival time of \( F' \) from \( t' \) to \( t_0 \), it still completes at time \( t_{\text{end}'} \) since there is at least one frame with priority that is higher for \( i \) is executed. Therefore, frame \( F \) still completes at time \( t_{\text{end}'} \).

Assuming that for each frame with priority that is higher than \( i \), we move its first arrival time after \( t_0 \) to \( t_0 \). In this case, the completion time of frame \( F \) might be delayed to \( t_{\text{end}'} \). Moreover, their succeeding frames are requested early with their minimum separations, and each frame is executed for a maximum execution time. The completion time of frame \( F \) is delayed further to time \( t_{\text{end}''} \). Since time \( t \) is one of critical instant candidates of \( T \), the response time of frame \( F \) is longer than the first situation.

When we assume that frame \( F \) with priority \( i \) is the \( k \)th frame of the arrival of \( T_i \) after \( t_0 \), the worst case response time of \( F \) is caused. If there exists a frame between \( F \) and \( F' \) with priority that is lower than \( i \), \( t_0 \) exists between \( t \) and \( t' \). \( t \) is one of critical instant candidates of \( T \).

Based on the new critical instant theorem, a frame in the task set is schedulable when critical instant of the frame is smaller than its deadline.

### 3.2 Maximum Interference Function (MIF)

In order to decide the feasibility of a set of multi-frame tasks effectively, a maximum interference function (MIF) is proposed in [4]. The MIF of a task is a function presenting the maximum time during which the execution of the task interferes with the execution of the lower priority tasks within the duration time interval \( t \).

The MIF \( M_i(t) \) of a multi-frame task \( T_i \) is represented as follows, where \( I_k^i \) is called the interference function (IF) and represents the maximum time during which \( T_i \) interferes with lower priority tasks during \( t \) after the arrival of the \( k \)th frame of \( T_i \),

\[
M_i(t) = \max_{0 \leq k \leq N_i-1} I_k^i(t)
\]

It can be transformed into the following form.

\[
I_k^i(t) = \sum_{h=k}^{k+J-1} c_{ij}d_{ijh}^N + \min\left( C_{ij} + \sum_{h=k}^{k+J-1} t_{ij}^N, t - \sum_{h=k}^{k+J-1} t_{ij}^N \right)
\]

Where \( J \) is the maximum integer that satisfies the following formula.

\[
\sum_{h=k}^{k+J-1} t_{ij}^N \leq t
\]

The MIF of a multi-frame task is the maximum value of each interference function (IF) of frames when interference frames are released simultaneously. Consider a task consisting of three frames: \( \tau_i = ((1, 4, P1, 4), (2, 4, P2, 4), (3, 4, P3, 4)) \), Fig. 4 illustrates the IFs of three kinds of conditions and the MIF of \( \tau_i \). The MIF of the multi-frame task \( \tau_i \) is the uppermost connected lines of the three IFs.

From the definition of the MIF, the necessary and sufficient condition for the schedulability of a multi-frame task \( T_i \) can be given as follows.

A multi-frame task \( T_i \) is schedulable under a GMF model if and only if:

\[
\exists \tau, 0 \leq t \leq D_i^f, \sum_{j=1}^{i-1} M_j(t) + C_i \leq t
\]

According to the above formula, we can obtain schedulable conditions for a task during time interval \( t \). Since the interference time caused by one higher priority multi-frame task can be calculated, the total interference time caused by one or more than higher priority multi-frame tasks can be intuitively calculated by arithmetically adding of them. In fact, for a set of multi-frame tasks that has one or more tasks, the result of the arithmetical addition of MIFs is longer than the actual interference time caused by all the higher priority tasks. Saturation summation has been applied to improve MIF for a task set.

### 3.3 Saturation Summation

For a set of tasks, the degree of schedulable of a frame is captured by total interference time requested by all higher priority frames within the duration of time interval \( t \). As described above, the time is calculated by arithmetical addition of MIFs. The MIF of \( \tau_{1} \) and \( \tau_{2} \) is shown in Fig. 5 (a) (i.e., \( M(t) = M_{\tau_{1}}(t) + M_{\tau_{2}}(t) \)). \( M_{\tau_{1}}(t) \) constitutes the actual interference time affected by the two multi-frame task \( \tau_{1} \) and \( \tau_{2} \). From Fig. 5, the value of \( M(t) \) is greater than that of \( M_{\tau_{1}}(t) \) during the two intervals of time \( (0, 4) \) and \( (8, 11) \), which means that \( M(t) \) is longer than the actual interference time in the same two intervals. In
fact, MIF is impossible to increase like $M(t)$ during the two intervals of time, because only one frame is executed during the intervals. For this reason, $M(t)$ must be converted into $M_j(t)$ by use of a saturation conversion.

**Definition 2 (Saturation Conversion):** Given a function $F(t)$ that depends on time with a nonnegative integer gradient, a new function $F_j(t)$ can be obtained on the basis of the following formulas. This is called a saturation conversion.

$$F_j(t) = t - \max_{0 \leq \tau \leq t} (\tau - F(\tau))$$

As for the definition of saturation conversion, the function that is converted from the arithmetical addition of MIFs by a process of saturation conversion can be described as saturation summation, which can be expressed as follows:

$$\sum_{j=1}^{i-1} M_j(t) = t - \max_{0 \leq \tau \leq t} \left( \tau - \sum_{j=1}^{i-1} M_j(\tau) \right)$$  \hspace{1cm} (2)

Figure 5 (b) illustrates saturation summation applied to the same set of tasks in Fig. 5 (a).

- In the interval of time $[0, 5]$, because $\sum_1^3 M_j(\tau) \geq \tau$,

$$\max_{0 \leq \tau \leq 4} \left( \tau - \sum_{j=1}^{i-1} M_j(\tau) \right) = 0, \quad \Rightarrow \sum_{j=1}^{i-1} M_j(t) = t; \quad (0 \leq t \leq 4)$$

- In the interval of time $[8, 11]$, because,

$$\max_{8 \leq \tau \leq 11} \left( \tau - \sum_{j=1}^{i-1} M_j(\tau) \right) = 3, \quad \Rightarrow \sum_{j=1}^{i-1} M_j(t) = t - 3; \quad (8 \leq t \leq 11)$$

Thus, saturation summation of MIF can compensate for a discrepancy and calculate interference time precisely.

The necessary and sufficient condition for the schedulability of a $k$th frame of a multi-frame task $T_i$ must be reformulated for saturation summation of MIF.

According to the definition of saturation summation, the inequality 1 can be transformed into the following inequality.

$$\exists t, 0 \leq t \leq D_i^k, \quad \sum_{j=1}^{i-1} M_j(t) + C_i^j \leq t$$  \hspace{1cm} (3)

Since saturation summation is a monotonic non-decreasing function, the laxity of a multi-frame task $T_i$ is at the maximum at a time $t$ is equal to the task deadline. Based on the above discussions, it is easy to observe that saturation summation of MIF can only be inspected at the time of the deadline $t$, but not at each $t' < t$. The necessary and sufficient condition for the schedulability of the $k$th frame of a multi-frame task $T_i$ is described below:

**Theorem 2:** The $k$th frame of a multi-frame task $T_i$ can be schedulable if and only if

$$\sum_{j=1}^{i-1} M_j(D_i^k) + C_i^j \leq D_i^k$$

This theorem is proved by demonstrating the two inequalities (1) and (3). When inequality (1) is satisfied, then inequality (3) is also satisfied. The two inequalities are equivalent to the schedule conditions of the $k$th frame of a multi-frame task $T_i$. The proof of this is described in details elsewhere [5].

**4. Scheduling Algorithms**

When priorities for a frame are assigned sequentially, the worst-case complexity of searching the frame deadline with $n$ frames is of a higher order than $O(n!)$. It took several days to calculate a 10 frames problem (4 I/O blocking tasks and 2 non-I/O blocking tasks). The algorithm is practically unusable for most real-time applications.

**4.1 Frame Laxity Monotonic Scheduling (FLMS)**

In this section, we present a minimum laxity first dynamic scheduling algorithm based on a greedy approach such as [14]. Laxity is a measure of the flexibility available for scheduling a task. The basic idea is that the shorter laxity of frame is, the higher priority assigned to the frame. In general, laxity $X_i$ of a non-I/O blocking $ith$ frame is the maximum time that the frame can be delayed on its activation to complete within its deadline. For an I/O blocking task
Algorithm: FindUrgencyFrame
Input: The minimum laxity frame \( f_{min} \)
Output: An urgency frame \( f \) if it exists;
1. For each \( i \)th unselected frame in the set of frames \( F \) \{
2. Assume the \( i \)th frame is in the \( i \)th task;
3. if \( \text{MIF}_{f_{min}} + \) the execution time of the \( i \)th frame \( \geq \) The deadline of the \( i \)th task \} 
4. The urgency frame \( f_u \) is the \( i \)th frame; 
5. Return \( f_u \); 
6. \}
7. \}

\[ X_f = \begin{cases} 
D_k - \text{WCIT}(f) - L_f - C_f - B(f) & \text{if } P_f' \text{ is always fixed.} \\
(D_k - \text{WCIT}(f) - C_f - C_f' - B_k)/2 & \text{if } P_f' \text{ and } P_f \text{ are not fixed.} 
\end{cases} \]

Where, \( D_k \) is a task deadline, \( \text{WCIT}(f) \) is amount of time during which interference is caused by the execution of higher priority frames in worst case, as calculated by means of theorem 1. \( L_f \) is the length of frame \( f' \). \( B(f) \) is a function, when frame \( f \) is the frame applicable before I/O blocking, \( B(f) \) returns a value of a maximum waiting time (\( B_k \)) of the task; otherwise, \( B(f) \) returns zero.

According to the above formula, if priorities of \( f' \) and \( f \) are not assigned, frame laxity values of \( f' \) and \( f \) are equal to the average equivalent of the laxity of the I/O blocking task. Note that, \( B_k \) is the maximum waiting time.

If more than one frame have the same laxity, the algorithm randomly selects the one to which higher priority should be given.

To consider the I/O blocking task mechanism, we improved FLMS based on the following strategies:

- FLMS takes into consideration the execution time of a frame. However it can not ensure that other frames do not miss deadlines. Consider two frames \( \tau_1 \) (1, 100, \( P_1 \), 100) and \( \tau_2 \) (52, 150, \( P_2 \), 150). The interference time caused by higher priority frames is 50. The laxity of \( \tau_1 \) and \( \tau_2 \) are 49 and 48, respectively. According to FLMS, \( \tau_2 \) is assigned a higher priority. But \( \tau_1 \) is unschedulable. This example illustrates that a frame with minimum laxity is not necessary an “urgency” frame that must be executed in the first instant, the “urgency” frame does not always exist. We compare the total of interference time and execution time of the frame with minimum laxity with the deadlines of other frames so as to find an “urgency” frame. If the deadline of a frame is smaller than the total value, the frame is deemed to be an “urgency” frame, and must be assigned higher priority. Otherwise, the frame with the minimum laxity is assigned the higher priority (line 12).

- The two frames that are in an I/O blocking task do not mutually interfere with each other regardless of the priority of the two frames. In other words, when calculating interference time of the \( i \)th frame, it is unnecessary to generate the IF of the multi-frame task to which the \( i \)th frame belongs (line 4).

- Whenever a frame fails to meet its deadline, the equal slack allocation proposed in [3] is used to maximize the schedulability of lower priority frames. If the priority of two frames of an I/O blocking task is greater than that of a frame that fails to meet its deadline, the deadline of the two frames is adjusted by this method (line 17).

According to FLMS mechanism, the minimum laxity frame in a set of frames in which frames are not assigned priorities is taken into account (from line 2 to line 7). If the frame is unschedulable, showing the whole frame set is unschedulable, the algorithm is stop at this step (line 9). Otherwise, finding an urgency frame is performed at line 11. If the urgency frame is existed, it is used to replace the minimum laxity frame. From line 11 to line 14, priority of the frame we selected is assigned and deadline of the frame is determined. At the end, the equal slack allocation is performed if the frame we selected is in an I/O blocking task (line 17). Note that priority of a frame is not equal to that of another frame, since the relevant MIF calculation becomes more complex. The method is summarized in Fig. 7.

4.2 Genetic Algorithm for I/O Blocking

Our GA requires the definition of a set of genetic operations and an evaluation function as follows:
4.2.1 Coding Individual

Each frame is denoted with frame identifier (FrameID). Based on the frame identifier, we can find whether an I/O task that the frame belongs to. Priorities of frames can be encoded by a string where the \( i \)th number in string is priority of the \( i \)th frame. Another string is used for encoding deadlines of frames. The lengths of two strings are the number of frames. The deadline of a frame with the highest priority is equal to the maximum execution time of the frame. Priorities and deadlines are generated randomly in their possible ranges. According to the FrameID, priority and deadline of the frame are easy to find.

4.2.2 Generate Initial Generation

Individuals are generated randomly as many as population size. The fitness value of an individual is calculated by an evaluation function. When the same individual is generated, it is deleted and a new different individual is obtained by using mutation operation. Then individuals in the generation will be sorted in ascending order.

4.2.3 Evaluation Strategy

The fitness value of an individual is the number of frames that are schedulable in the task set. The schedulability of a frame is calculated according to the new critical instant theorem.

4.2.4 Crossover

The crossover operation is held between two parents. Two selected individuals are called parent-A and parent-B, and assume the fitness value of parent-A is higher or equal to that of parent-B. Every crossover results in one offspring by randomly choosing one of two suboperations. One is to exchange priorities of some genes. This mechanism can be carried out by cyclic permutation operation. The same performance was described elsewhere [15] in details.

Another operation is to change deadlines of some genes by the following steps:

1. For a gene \( d_a \) in parent-A and the gene \( d_b \) at the same position as \( d_a \) in parent-B, if \( d_a = d_b \) then set \( d_a \) to the same position in the offspring.
2. Let the number of unset genes in the offspring be \( m \).
   
   (a) if \( m \leq 4 \) then for every unset positions in the offspring, set the average value of genes at the same position in parents to the offspring.
   
   (b) if \( m > 4 \) then: select \( [m + 1]/4 \) genes from parent-A randomly and set them to the offspring. If the deadline of a frame in an I/O blocking task was determined, the deadline of another frame in the same I/O blocking task is also determined. Values of unset genes in offspring are copied from genes at the same position in parent-B.

At the end of crossover operation, let the deadline of a frame with the highest priority equal to the maximum execution time of the frame.

4.2.5 Mutation

The mutation operation is in an individual. We also conduct mutation operation by randomly choosing one of two suboperations. One is to choose an I/O blocking task randomly, change deadlines of two frames randomly. Another is to choose an I/O blocking task randomly, set the average value of deadlines in two frames.

4.2.6 Other Operations

Selection and differentiation of the same individual are performed as described elsewhere [16] in details.

Parameters for GA were set as follows. Population size was five times as much as the number of frames. The probability of mutation was 0.2. The number of generations was 1000.

5. Experiments

This section presents a set of experiments demonstrating the effectiveness of our approach. The two developed algorithms are implemented using C language executed on a 2.53 GHz Pentium4 processor with a 512 MB RAM. We have not conducted a comparison with other algorithms. Because I/O blocking with a multi-frame task model is different from L&L task model and two algorithms we propose are also based on the new critical instant theorem. We conducted tests on four task sets consisting of 5, 10, 15 and 20 frames, respectively. 20 samples were generated randomly for each task set. The main parameters of a task are listed in Table 1. The minimum execution time and the maximum execution time of frame are 10 and 100 time units, respectively.

The number of frames, I/O blocking tasks and non-I/O blocking tasks in the four task sets are listed in Table 1. The task deadlines and the blocking waiting times were generated randomly between the minimum time and the maximum time of the deadline and the waiting time, respectively.

The data reported in Figs. 8 and 9 are the average values of 20 samples obtained by FLMS and GA for each data set.

<table>
<thead>
<tr>
<th>Name of data sets</th>
<th>Set 1</th>
<th>Set 2</th>
<th>Set 3</th>
<th>Set 4</th>
</tr>
</thead>
<tbody>
<tr>
<td>Number of frames</td>
<td>5</td>
<td>10</td>
<td>15</td>
<td>20</td>
</tr>
<tr>
<td>Number of I/O blocking tasks</td>
<td>2</td>
<td>4</td>
<td>7</td>
<td>9</td>
</tr>
<tr>
<td>Number of non-I/O blocking tasks</td>
<td>1</td>
<td>2</td>
<td>1</td>
<td>2</td>
</tr>
<tr>
<td>Minimum deadline for task</td>
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<td>400</td>
<td>700</td>
<td>1300</td>
</tr>
<tr>
<td>Maximum deadline for task</td>
<td>500</td>
<td>1000</td>
<td>1300</td>
<td>2000</td>
</tr>
<tr>
<td>Minimum waiting time</td>
<td>20</td>
<td>20</td>
<td>20</td>
<td>80</td>
</tr>
<tr>
<td>Maximum waiting time</td>
<td>90</td>
<td>90</td>
<td>90</td>
<td>150</td>
</tr>
</tbody>
</table>
In this paper, effective algorithms that can schedule a set of tasks which has I/O blocking under dynamic priority assignment are proposed. We have presented a new critical instant theorem necessary for a multi-frame task set with arbitrary priority assignment. An approach to schedulability analysis based on the critical instant theorem is also presented. In addition, we have proposed an FLMS algorithm. A genetic algorithm is also applied. From our experiments, we can conclude that the FLMS can significantly reduce the time of the calculation time, and GA can improve task schedulability ratio more than is possible with FLMS.

6. Conclusions

In this paper, effective algorithms that can schedule a set of tasks which has I/O blocking under dynamic priority assignment are proposed. We have presented a new critical instant theorem necessary for a multi-frame task set with arbitrary priority assignment. An approach to schedulability analysis based on the critical instant theorem is also presented. In addition, we have proposed an FLMS algorithm. A genetic algorithm is also applied. From our experiments, we can conclude that the FLMS can significantly reduce the time of the calculation time, and GA can improve task schedulability ratio more than is possible with FLMS.

References

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